

changes to RHS:

Consider

$$\begin{array}{ll} \text{Min} & c^T x \\ \text{st.} & Ax = b \\ & x \geq 0 \end{array} \quad (P)$$

$$\begin{array}{ll} \text{Min} & c^T x \\ \text{st.} & Ax = b + \theta \cdot e_i \\ & x \geq 0 \end{array} \quad (P')$$

Theorem: Let B be optimal for (P) .
 B is optimal for (P') if and only if

$$\underbrace{B^{-1}b}_{\bar{b}} + \theta \underbrace{B^{-1}e_i}_{i^{\text{th}} \text{ column of } B^{-1}} \geq 0$$

[From lecture 12']

proof. B is optimal for (P)

$$\Rightarrow \bar{b} = B^{-1}b \geq 0 \quad (1)$$

$$\bar{c}^T = c^T - c_B^T B^{-1}A \geq 0 \quad (2)$$

B is optimal for (P') iff

$$\begin{aligned} \bar{b}' &= B^{-1}b' = B^{-1}(b + \theta e_i) \\ &= B^{-1}b + \theta B^{-1}e_i \geq 0 \end{aligned}$$

$$\begin{aligned} \bar{c}' &= c'^T - c_B'^T B^{-1}A = c^T - c_B^T B^{-1}A \\ &= \bar{c} \geq 0 \end{aligned}$$

always holds, by (2)

[From lecture 12]

Observe that even if B stays optimal (P'),

$$\bar{x}'_B = B^{-1}b' = B^{-1}(b + \theta e_i) = \underbrace{B^{-1}b}_{\bar{x}_B} + \theta B^{-1}e_i \neq \bar{x}_B$$

However,

$$\bar{y}' = (B^T)^{-1}c'_B = (B^T)^{-1}c_B = \bar{y}$$

By strong duality $z' = c'^T \bar{x}' = b'^T \bar{y}'$

$$= (b + \theta e_i)^T \bar{y}$$

$$= b^T \bar{y} + \theta e_i^T \bar{y}$$

$$= z + \theta \bar{y}_i$$

[from lecture 12]

EXAMPLE: House Depot produces:

- hammers (fixed price 130)
from 1.5 kg of steel, 1 rivet, 0.3 kg of plastic, and
- pliers (fixed price 100)
from 1 kg of steel, 1 rivet, 0.5 kg of plastic.

Current stocks are:

27 kg of steel, 21 rivets, 9 kg of plastic.

Maximize income from current stocks.

(Assume that fractional hammers and pliers are ok)

[From lecture 11]

VAR:
 x_1 : hammers
 x_2 : pliers

MODEL:
MAX 130 · x_1 + 100 · x_2
S.t. 1.5 x_1 + 1 x_2 ≤ 27 ← steel
1 x_1 + 1 x_2 ≤ 21 ← rivets
0.3 x_1 + 0.5 x_2 ≤ 9 ← plastic
 x_1 , x_2 ≥ 0

[From lecture 11]

Example: hammers & pliers (cont'd)

S.E.F.:

$$\text{Min } [-130 \quad -100 \quad 0 \quad 0 \quad 0] x$$

$$\text{s.t. } \begin{bmatrix} 1.5 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0.3 & 0.5 & 0 & 0 & 1 \end{bmatrix} x = \begin{bmatrix} 27 \\ 21 \\ 9 \end{bmatrix}$$

$x \geq 0$

[simplex method \Rightarrow]

optimal basis: $B = \{1, 2, 5\}$

$$x^* = (12, 9, 0, 0, 0.9)$$

$$y^* = (-60, -40, 0)$$

$$z^* = -2460 \leftarrow \text{-profit}$$

optimal tableau:

$$\text{Min } [0 \quad 0 \quad 60 \quad 40 \quad 0]x$$

$$\begin{bmatrix} 1 & 0 & 2 & -2 & 0 \\ 0 & 1 & -2 & 3 & 0 \\ 0 & 0 & 0.4 & -0.9 & 1 \end{bmatrix} x = \begin{bmatrix} 12 \\ 9 \\ 0.9 \end{bmatrix}$$

$$x \geq 0$$

optimal basis matrix:

$$B = \begin{bmatrix} 1.5 & 1 & 0 \\ 1 & 1 & 0 \\ 0.3 & 0.5 & 1 \end{bmatrix}$$

$$B^{-1} = \begin{bmatrix} 2 & -2 & 0 \\ -2 & 3 & 0 \\ 0.4 & -0.9 & 1 \end{bmatrix}$$

Q1: We change the available steel from 27 to $27 + \theta$.

For what values of θ does the basis remain optimal?

$$b' = b + \theta \cdot e_1$$

$$= \begin{bmatrix} 27 \\ 21 \\ 9 \end{bmatrix} + \theta \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

$$\bar{b}' = B^{-1} \cdot b' = B^{-1} \cdot (b + \theta e_1)$$

$$= \underbrace{B^{-1}b}_b + \theta \underbrace{B^{-1}e_1}_{\text{first column of } B^{-1}}$$

$$= \begin{bmatrix} 12 \\ 9 \\ 0.9 \end{bmatrix} + \theta \begin{bmatrix} 2 \\ -2 \\ 0.4 \end{bmatrix} \geq 0$$

$$\begin{cases} 12 + \theta \cdot 2 \geq 0 \\ 9 + \theta \cdot (-2) \geq 0 \\ 0.9 + \theta \cdot 0.4 \geq 0 \end{cases}$$

implied by
(3)

$$\Leftrightarrow \begin{cases} \theta \geq -6 \\ \theta \leq 4.5 \\ \theta \geq -2.25 \end{cases}$$

↑
(1)
(2)
(3)

$$\Leftrightarrow \theta \in [-2.25, 4.5]$$

Definition: The set of all θ such that B stays optimal is called the allowable range.

Q2. What is the allowable range for plastic ?

$$b' = \begin{bmatrix} 27 \\ 21 \\ 9 \end{bmatrix} + \theta \begin{bmatrix} 0 \\ 0 \\ -1 \end{bmatrix} \leftarrow e_3$$

$$\bar{b}' = \begin{bmatrix} 12 \\ 9 \\ 0.9 \end{bmatrix} + \theta \begin{bmatrix} 0 \\ 0 \\ -1 \end{bmatrix} \geq 0 \leftarrow \text{Third column of } B^{-1}$$

$$\begin{cases} 12 + 0 \cdot \theta \geq 0 \\ 9 + 0 \cdot \theta \geq 0 \\ 0.9 + 1 \cdot \theta \geq 0 \end{cases} \Leftrightarrow$$

$$\begin{cases} 12 \geq 0 \\ 9 \geq 0 \\ \theta \geq -0.9 \end{cases} \text{ always true}$$

$$\theta \in [-0.9, +\infty[$$

Q3. A seller proposes 4 additional kg of steel for \$200. Should you accept?

Note: from a_1 , $\theta = 4 \in [-2.25, 4.5]$,
i.e. 4 is in allowable range for steel.

$$\begin{aligned} z^{1*} &= [-130 \quad -100 \quad 0 \quad 0 \quad 0] \cdot x^* = b^{1T} \cdot y^{1*} \\ &= (b + \theta \cdot e_1)^T y^* \\ &= b^T y^* + \theta e_1^T y^* \\ &= z^* + \theta \cdot y_1^* \end{aligned}$$

$$= -2460 - 60. \text{€}$$

$$= -2460 - 240$$

(-profit) decreases by
240

profit increases by
240

→ ACCEPT!

Q4

Someone is willing to buy 2 kg of steel off our stock for \$130.

Accept?

Note: $\theta = -2 \in [-2.25, 4.5]$

→ basis stays the same

$$\begin{aligned} z'^* &= z^* + \theta \cdot y_i^* \\ &= -2460 + (-2) \cdot (-60) \\ &= -2460 + 120 \end{aligned}$$

(-profit) increases by 120
profit decreases by 120 < 130
→ ACCEPT!

changes to objective function

Consider

$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & Ax = b \\ & x \geq 0 \end{aligned} \quad (P)$$

$$\begin{aligned} \min \quad & (c + \theta \cdot e_j)^T x \\ & Ax = b \\ & x \geq 0 \end{aligned} \quad (P')$$

for $\theta \in \mathbb{R}$

Theorem Let B be optimal for (P) .

B is optimal for (P') if and only if

$$\left\{ \begin{array}{l} \bar{c}_j + \theta \geq 0 \quad \text{if } j \notin B \\ \bar{c}_N - \theta \underbrace{e_i^T B^{-1} N}_{\text{if } j \text{ is the } i^{\text{th}} \text{ basic column.}} \geq 0 \end{array} \right.$$

nonbasic part of i^{th} row of optimal tableau.

Proof

B is optimal for (P) \Rightarrow

$$\bar{b} = B^{-1}b \geq 0 \quad (1)$$

$$\bar{c}^T = c^T - c_B^T B^{-1}A \geq 0 \quad (2)$$

B is optimal for (P') iff

$$\bar{b}' = B^{-1}b' = B^{-1}b = \bar{b} \geq 0$$

always holds, by (1)

$$\bar{c}' = c'^T - c_B'^T B^{-1}A \geq 0$$

Case 1: $j \notin B$

$$\bar{c}' = (c + \theta e_j)^T - c_B^T B^{-1} A$$

$$= \underbrace{c^T - c_B^T B^{-1} A}_{\bar{c}} + \theta e_j^T$$

$$= \underbrace{\bar{c}^T}_{\geq 0} + \theta e_j^T \geq 0$$

$$\Leftrightarrow \bar{c}_j + \theta \geq 0$$

Case 2: $j \in \mathcal{B}$, j is the i^{th} basic column.

$$\begin{aligned}\bar{c}' &= (c + \theta e_j)^T - (c_{\mathcal{B}} + \theta e_i)^T B^{-1}A \\ &= \underbrace{c^T - c_{\mathcal{B}}^T B^{-1}A}_{\bar{c}} + \theta e_j^T - \theta e_i^T B^{-1}A\end{aligned}$$

$$= \bar{c}^T - \theta \left(\underbrace{e_i^T B^{-1}A}_{\substack{i^{\text{th}} \text{ row} \\ \text{of optimal tableau}}} - e_j^T \right) \geq 0$$

i^{th} row of optimal tableau

i^{th} row of optimal tableau with a zero for basic variable

$$\Leftrightarrow \bar{c}_N^T - \theta \cdot e_i^T B^{-1}N \geq 0$$

Q5

What is the allowable range for c_1 ?

Note: $1 \in B$, first basic variable

$$\begin{aligned}\bar{c}_1^T &= c_1^T - c_B^T B^{-1}A \\ &= (c_1^T + \theta e_1^T) - (c_B^T + \theta e_1^T) B^{-1}A \\ &= c_1^T - c_B^T B^{-1}A - \theta (e_1^T B^{-1}A - e_1) \\ &= \bar{c}_1^T - \theta ([\cancel{x} \ 0 \ 2 \ -2 \ 0] - [\cancel{x} \ 0 \ 0 \ 0 \ 0]) \\ &= [0 \ 0 \ 60 \ 40 \ 0] - \theta [0 \ 0 \ 2 \ -2 \ 0] \\ &\geq 0\end{aligned}$$

$$\Leftrightarrow \begin{cases} 0 - 0.\theta & \geq 0 \\ 0 - 0.\theta & \geq 0 \\ 60 - 2.\theta & \geq 0 \\ 40 + 2.\theta & \geq 0 \\ 0 - 0.\theta & \geq 0 \end{cases}$$

basic columns
trivially true

$$\Leftrightarrow \begin{cases} \theta \leq 30 \\ \theta \geq -20 \end{cases} \Leftrightarrow \theta \in [-20, 30]$$

Warning:

in terms of original costs
(max), we have

$$-\theta \in [-30, 20]$$